# On some Graph Colo(u)ring Problems 

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## Quick Overview - Acyclic Coloring

## Definition

Acyclic coloring of a graph $G$ :

- Proper vertex coloring of $G$ (no two neighbors are assigned the same color)
- Any two classes of colors induce an acyclic graph (=a forest)
- Equivalently, no bicolored cycle in $G$
- $a(G)$ : smallest number of colors necessary to acyclically color $G$


## Acyclic Coloring

- (Asymptotically tight) bounds on $a(G)$ for $G=d$-dimensional grid, $d \geq 2$
- $a(G) \leq 9$ for any graph of max. degree 5 (recently improved to 8 )


## Quick Overview - Star Coloring

## Definition

Star coloring of a graph $G$ :

- Proper vertex coloring of $G$
- Any two classes of colors induce a forest of stars
- Equivalently, no bicolored path of length 3 in $G$
- $\chi_{s}(G)$ : smallest number of colors necessary to star color $G$


## Star Coloring

- Lovász Local Lemma $\Rightarrow$ general upper bound on $\chi_{s}$ for graphs of max degree $\Delta$
- $\chi_{s}(G)=O\left(\Delta^{\frac{3}{2}}\right)$
- Lower bound: $\chi_{s}(G)=\Omega\left(\frac{\Delta^{\frac{3}{2}}}{\log \Delta^{\frac{1}{2}}}\right)$


## Quick Overview - Oriented Coloring

## Definition

Oriented coloring of a graph $G$ :

- Let $\vec{G}$ be an orientation of $G$
- Proper vertex coloring of $\vec{G}$
- If $(u, v),(x, y) \in A(\vec{G})$, we cannot have $c(u)=c(y)$ and $c(v)=c(x)$
- $\chi_{o}(G)$ : smallest number of colors necessary to color any orientation $\vec{G}$


## Oriented Coloring

- Bounds for the oriented coloring of 2-D grids:

$$
7 \leq \chi_{0}\left(G\left(n_{1}, n_{2}\right)\right) \leq 11
$$

- Conjecture: $\chi_{o}\left(G\left(n_{1}, n_{2}\right)\right)=7$ by homomorphism to $T(7 ; 1,2,4)$


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## Quick Overview - $L(p, q)$-labeling

## Definition

$L(p, q)$-labeling of a graph $G$ :

- Vertex labeling of $G$, in the (integer) range $[0 ; M]$
- Neighbors in $G$ must be assigned values that differ by at least $p$
- Vertices at distance 2 in $G$ must be assigned values that differ by at least $q$
- $\lambda_{q}^{p}(G)$ : smallest value of $M$ necessary to $L(p, q)$-label $G$


## Quick Overview - $L(p, q)$-labeling

## $L(p, q)$-labeling

| Resp. values of $p$ and $q$ | $\lambda_{q}^{p}\left(G_{d}\right) \geq$ | $\lambda_{q}^{p}\left(G_{d}\right) \leq$ | Gap | No-hole Property |
| :---: | :---: | :---: | :---: | :---: |
|  | $p$ |  | 0 | No (except for $p=1$ ) |
| $p=0 ; q \neq 0$ | $(2 d-1) q$ |  | 0 | No (except for $q=1$ ) |
| $p, q \geq 1 ; 2 p<q$ | $2 p+(2 d-2) q$ | $2 d q$ | $2 q-2 p$ | No |
| $q=1 ; 1 \leq p \leq 2 d$ | $2 p+2 d-2$ |  | 0 | Yes (Prop 4) |
| $p, q \geq 1 ; q \leq 2 p \leq 4 d q$ | $2 p+(2 d-2) q$ | $2 p+(2 d-1) q-1$ | $q-1$ | Yes (Prop 4) |
| $p, q \geq 1 ; p=\alpha q ; \alpha \leq 2 d$ | $2 p+(2 d-2) q$ |  | 0 | No |
| $\begin{aligned} & p, q \geq 1 ; p=\alpha q+\beta \\ & 1 \leq \beta \leq q-1 ; p \leq 2 d q+\beta-q \end{aligned}$ | $2 p+(2 d-2) q$ | $2 p+2 d q-2 \beta$ | $2(q-\beta)$ | No |
| p,q $\geq 1 ; p \geq 2 d q+1$ | $p+(4 d-2) q$ |  | 0 | No |

Table 1: $L(p, q)$ labeling of $G_{d}$ : Summary of the results

## The S-labeling problem (joint work with S. Vialette)

## Notations

- Let $G$ be a graph having $n$ vertices and $m$ edges
- An $S$-labeling $\theta: V(G) \rightarrow\{1,2 \ldots n\}$ is a bijective mapping of the vertices of $G$
- Let $\Theta(G)$ be the set of all $S$-labelings of $G$
- For any graph $G$ and any $\theta \in \Theta(G)$,

$$
\mathrm{SL}(G, \theta)=\sum_{e \in E(G)} \min \{\theta(u): u \in e\}
$$

## The S-Labeling Problem

Given a graph $G$, find an $S$-labeling $\theta \in \Theta(G)$, such that $\operatorname{SL}(G, \theta)$ is minimized.

## S-labeling of $K_{1,6}$

## Example


$1+2+3+4+4+4=18$

$1+1+1+1+1+1=6$

Figure: 2 different $S$-labelings - One of them is optimal

## Overview of the Results

## Theorem

The S-labeling problem is:

- NP-complete for general graphs [Vialette 06]
- polynomial-time solvable for split graphs


## Overview of the Results

## Definition

- Split graph = chordal graph whose complement is also chordal
- Vertices of a split graph can be partitioned into:
- A clique $K$, and
- A stable set $S$


## Example



Figure: A split graph $G$ and a partition $V(G)=K \cup S$

## Overview of the Results

## Theorem

For any graph $G$ of maximum degree $\Delta$,

$$
\frac{m(m+\Delta)}{2 \Delta} \leq S L(G) \leq \frac{m(n+1)}{3}
$$

## Theorem

Approximation ratios for the S-labeling problem:

- $r=\frac{4}{3}$ for general graphs
- $r=\frac{2 \Delta}{3}$ for trees of max degree $\Delta$
- $r=\frac{4 \Delta}{3}$ for graphs of max degree $\Delta$


## Overview of the Results

## Definition (Contribution Vector)

For any $S$-labeling $\theta$ of a graph $G$, and for any $1 \leq i \leq|V(G)|, C V_{\theta}[i]=$ number of times label $i$ contributes to the sum in $\operatorname{SL}(G, \theta)$.

## Example



Figure: $C V_{\theta}=[2,2,2,2]$ (left) ; $C V_{\theta}=[4,1,1,1,1]$ (right)

## Overview of the Results

## Theorem

For any graph $G$ and any optimal S-labeling $\theta$ of $G$, its contribution vector $\mathrm{CV}_{\theta}$ is monotone non-increasing

## Theorem

For any graph $G$ of maximum degree $\Delta$, any optimal S-labeling $\theta$ of $G$ satisfies $C V_{\theta}[1]=\Delta$

## Conclusion on the S-labeling problem

## Some open questions

- Let $k$ be s.t.
- $C V_{\theta}[1]=C V_{\theta}[2]=\ldots=C V_{\theta}[k]=\Delta$, and
- $C V_{\theta}[k+1] \neq \Delta$
- Let $G_{\Delta}$ be the subgraph of $G$ induced by the vertices of degree $\Delta$
- Conjecture: there exists an optimal $S$-labeling $\theta$ s.t. the vertices associated to $C V_{\theta}[1 \ldots k]$ form a maximum independent set in $G_{\Delta}$
- Question 1: complexity of the $S$-labeling problem for trees ?
- Question 2: does the problem admit a PTAS ? a constant approximation ratio ?


## Conclusion

## Summary

- Glad and honored to pay tribute to André
- He's the one who taught me research
- Helped me at the beginning of my career
- Relations evolved from PhD supervisor, to colleague and friend


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## Thank you, André

